

1 Solution: produce or learn?

The **optimal strategy** has a very simple form: there is a time $t^* \in \{1, \dots, T\}$ such that (* denotes optimal quantities):

$$\begin{aligned} \text{before time } t^* & : \text{ only learn, i.e., } l^*(t) = 1 \text{ for } t < t^* \\ \text{from time } t^* & : \text{ only produce, i.e., } p^*(t) = 1 \text{ for } t \geq t^* \end{aligned}$$

This start-production time t^* can be explicitly determined (see proof below). Except possibly at time t^* , it is never optimal to split time between learning and producing ($l^*(t), p^*(t) \in (0, 1)$).

The optimal strategy has the following interpretation:

$$\begin{aligned} \text{learn at full speed } (l^*(t) = 1) & \quad \text{if} \quad \text{present quality is less than future capacity} \\ \text{produce at full speed } (p^*(t) = 1) & \quad \text{if} \quad \text{present quality exceeds future capacity} \end{aligned}$$

Let's first explain this interpretation more precisely and then prove that this strategy is optimal.

1.1 Interpretation

Let's call $Q(t) := P(t) + mL(t)$ the *quality at time t*. Since the objective function is strictly increasing in $p(t)$ and $l(t)$, at optimality, $p^*(t) + l^*(t) = 1$. We can hence eliminate $l(t)$ by substituting $l(t) = 1 - p(t)$ into $Q(t)$ to get:

$$Q(t) = \sum_{s=1}^{t-1} p(s) + m \sum_{s=1}^{t-1} (1 - p(s)) = m(t-1) - (m-1) \sum_{s \leq t-1} p(s)$$

Note that $Q(1) = P(1) + mL(1) = 0$ since $p(0) := 0 =: l(0)$. Our problem is then reduced to:

$$\begin{aligned} \max_{p(t), t=1, \dots, T} \quad & f(p) := \sum_{t=1}^T p(t) Q(t) \\ \text{subject to} \quad & 0 \leq p(t) \leq 1, \quad t = 1, \dots, T \end{aligned}$$

Taking partial derivatives of the objective function:

$$\frac{\partial f}{\partial p(t)} = Q(t) - (m-1) \sum_{s \geq t+1} p(s) =: Q(t) - \bar{Q}(t) \quad (1)$$

where the *marginal capacity-to-go* (from time $t+1$ to T) is defined as

$$\bar{Q}(t) := (m-1) \sum_{s \geq t+1} p(s)$$

If the quality $Q(t)$ at the present time t is less than the marginal capacity-to-go $\bar{Q}(t)$, i.e., if $Q(t) < \bar{Q}(t)$, then

$$\frac{\partial f}{\partial p(t)} < 0$$

and hence a smaller $p(t)$ will increase the objective value f . This means the optimal production effort $p^*(t) = 0$, i.e., learn fast and build up quality before producing. Conversely if $Q(t) > \bar{Q}(t)$, then $\partial f / \partial p(t) > 0$ and $p^*(t) = 1$, i.e., produce at full speed if present quality exceeds the marginal capacity-to-go.

1.2 Proof

To solve for the optimal $p^*(t)$ more explicitly, we note from (1) that

$$\frac{\partial f}{\partial p(t)} = m(t-1) - (m-1)(P - p(t)) \quad (2)$$

where $P := \sum_{t=1}^T p(t)$ is the *total production capability*. Our solution follows from the following properties.

It is optimal to initially learn and don't produce. In that case, the marginal production value increases rapidly at rate m , until eventually produce at full speed. More precisely, we have

Proposition 1 *Suppose $T \geq 2$. Then, at optimality,*

1. $p^*(1) = 0$.
2. If $p^*(t) = 0$ then

$$\frac{\partial f}{\partial p(t+1)} - \frac{\partial f}{\partial p(t)} \geq m$$

3. There exists some time $t \leq T$ at which $p^*(t) = 1$.

Proof. We have from (2), at optimality,

$$\frac{\partial f}{\partial p(1)} = -(m-1)(P^* - p^*(1)) = -(m-1) \sum_{s=2}^T p^*(s) < 0$$

where the last inequality holds as long as $T \geq 2$. Hence $p^*(1) = 0$.

The second claim follows directly from (2). To prove the third claim, suppose $p^*(s) = 0$ for $s = 1, 2, \dots, t-1$. Then $P^* - p^*(t) = \sum_{s=t+1}^T p^*(s)$. We have from (2)

$$\frac{\partial f}{\partial p(t)} = m(t-1) - (m-1) \sum_{s=t+1}^T p^*(s)$$

Hence if $p^*(t) = 0$ for all $t = 1, \dots, T$, then $\frac{\partial f}{\partial p(t)}(T) = m(T-1) > 0$. This means $p^*(T) = 1$, a contradiction. Hence there exists an $t \leq T$ such that $p^*(t) = 1$. ■

Proposition 1 says that there is a time t^* at which the optimal strategy $p^*(t)$ switches from 0 to 1. What are the optimal values of $p^*(t)$ after time t^* ? We claim that $p^*(t) = 1$ for all $t \geq t^*$.

Proposition 2 *Suppose $T \geq 2$. If $p^*(t) = 0$ for $t = 1, \dots, \tau - 1$ and $p^*(\tau) = 1$ for some τ , then $p^*(t) = 1$ for $t = \tau, \dots, T$.*

Proof. Since $p^*(t) = 0$ for $t = 1, \dots, \tau - 1$ and $p^*(\tau) = 1$, we have from (2), for $t = \tau, \dots, T$,

$$\frac{\partial f}{\partial p(t)} = m(t-1) - (m-1) \left(\sum_{s=\tau}^T p(s) - p^*(t) \right)$$

Compare this quantity at time τ and at any time $t > \tau$:

$$\frac{\partial f}{\partial p(t)} - \frac{\partial f}{\partial p(\tau)} = m(t-\tau) - (m-1)(p^*(\tau) - p^*(t)) > 0$$

where the last inequality follows because $t - \tau \geq 1$ but $p^*(\tau) - p^*(t) \leq 1$. Hence (noting that $p^*(\tau) = 1$)

$$\frac{\partial f}{\partial p(t)} > \frac{\partial f}{\partial p(\tau)} \geq 0$$

This implies $p^*(t) = 1$ for $t = \tau, \dots, T$. ■

The only remaining task is to determine the time t^* at which $p^*(t)$ switches from 0 to 1. Even though we don't know the value of $p^*(t^*)$, Propositions 1 and 2 imply that the values of $p^*(t)$ at other times are either 0 or 1:

$$p^*(t) = \begin{cases} 0 & \text{for } t = 1, \dots, t^* - 1 \\ 1 & \text{for } t = t^* + 1, \dots, T \end{cases}$$

Hence $P^* - p^*(t^*) := \sum_{s=1}^T p^*(s) - p^*(t^*) = T - t^*$. Hence, from (2), we have

$$\frac{\partial f}{\partial p(t^*)} = m(t^* - 1) - (m-1)(T - t^*) \quad (3)$$

The time t^* is when $\frac{\partial f}{\partial p(t)}$ switches from negative to positive. Since t^* must be an integer, (3) may not be 0. Let s^* be the real number that satisfies $m(s^* - 1) - (m - 1)(T - s^*) = 0$, i.e.,

$$s^* := \frac{(T + 1)m - T}{2m - 1} \quad (4)$$

Then $t^* = \lceil s^* \rceil$.

This completes the derivation of the optimal strategy.

Remark. How does the quality parameter m impact the start time t^* ? From (4), we see that, as m changes from 1 to ∞ , t^* changes from 1 to $(T + 1)/2$. Hence, as quality becomes more important (larger m), one starts to produce later in order to build up quality (up to roughly half the horizon T).